# Modified Sequence Algorithm for Computation of Optimum sequencing for Production scheduling

Naga Sai Ram.G, V.S. Haswanth.G, Durga Rajesh.K

**Abstract**— This paper focuses on the computation of the optimal sequence of the 2-Machines and n-Jobs. Conventional Johnson algorithm is long and often required time to compute them. Scaling such algorithm for m-machines n-Jobs is a bit complicated. The application of modified sequence algorithm to the n Jobs and m Machines is also made, which is the checking criteria for the n x3, n x5 and n xm scheduling problems change in the conventional algorithm by verifying the jobs value weather it is small and larger and appending them in priority sequence can help to solve such problems.

Index Terms— Optimum sequence, m-machines n-Jobs, Seheduling problem.

## 1 Introduction

Scheduling algorithms are used to find out the optimum sequence of the machines to be utilized for the optimum production. The application of the algorithm depends upon the number of jobs and machines in the problem. The basic possible case is n jobs 2 Machines where 2 machines are assigned by n jobs. Johnson algorithm is one such solving procedure for attaining the optimum solution of those problems. But this algorithm cannot be applied to the n jobs and m machines. The computed time for the Johnson algorithm depended on the number of searches taken to find the smallest time of in the job. In the computation terms this is represented by the  $O(N^2)$ . Which take time to find the optimum sequence for larger scaled problems. In order to decrease the time and the ease of computations of the sequence a small change has been made and rest of the algorithm is made to be same.

## **2** LITERATURE REVIRE

The first approach for the scheduling the two and three jobs production problem have been made by the Johnson S.M in 1954<sup>[1]</sup>. Solving the n jobs and m machines procedure have been given by the Herbert Campbell and Richard A.Duke without use of computer was published in 1970<sup>[2]</sup>.

#### **3 METHODOLOGY**

The basic work flow of this algorithm is to sort out the least time possible for the machines according to the jobs given. Which is the similar step in the normal Johnson algorithm. But the way of selecting the job is different in this method by selecting the smallest possible machine time in the machine cells and depending upon the positions of the smallest possible time appending the job with respect to the smallest entry. When it comes to the large scale problems it is easy to check the simple logic and decide the optimum sequence of the jobs. By using the divide and rule procedure the modified Johnson algorithm can be scaled up to m machines by using the resultant sequence is compared to the next machines and the procedure continues until the total number of the columns are converged. This algorithm can only be applied to the problem based on the time criteria.

Checking the optimal solution for the sequencing problem or deciding the optimal sequence based on the process cannot be the deciding factor, it is needed to check the all the possible cases say 2n where n is the total number of the jobs. This adaptation is not possible for the higher orders of n, because there lie only one optimum solutions in  ${}^{n}C_{1}$  where n is the number of the jobs.

sometimes chances of failing happen when made with conventional algorithm and alternate procedures needed to verify this cases and they may find much more optimal solution. So Modified Sequencing algorithm will be the checking criteria.

## **3.A) MODIFIED SEQUENCE ALGORITHM**

The modified sequence algorithms have the following operational steps:

- 1. Calculate number of columns which is *m*.
- 2. Calculate convergence limit which is *m*-1.
- 3. For *i*=0 compare *i*+1, *i*+2 columns for minimum row value.
- 4. If the minimum row value is in *i*+1 columns place the job value and respective minimum row value at front.
- 5. If the minimum row value is in *i*+2 place the job value and the respective minimum row value at end.
- 6. If the values of *i*+1 and *i*+2 are equal, then place the job value and respective row value at end.
- 7. Increment *i*.
- 8. Compare the newly formed list with *i*+*m*-1.
- 9. Continue step :3 until convergenc limit
- 10. Write only the job list.

we can apply the above algorithm only to the 2 Machines and n jobs which return the optimum sequence and this algorithm can be scaled to the n jobs and m machines.

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<sup>(</sup>This information is optional; change it according to your need.)

## **3.B) ANALYSIS OF THE ALGORITHM**

The time taken to solve the n x m job machine sequencing problems is very high both computationally and manually, whereas the modified sequencing algorithms takes very low time in either way. To prove the computational time of the modified sequencing algorithm the following considerations are taken.

For nxm job machine sequence matrix for  $T_c = [n+m-1]O(m)$ . Number of job interchange iterations is = n.

Number of comparison iterations is = m-1.

Time taken to process single line of iteration is O(m).

The total time complexity of the modified sequencing algrithm is

 $T_c = [n+m-1]O(m)$  ..... (2) where  $(m, n \ge 2)$ .

[n+m-1] is also called as the trans variable coefficient. The for the values of above and equal to 2 the progression starts as 3,5,7.....n<sup>th</sup> value is n+2 so the limiting factor of this algorithm is below 2. The modified sequencing algorithm is a linear algorithm. The advantages of this algorithm is as follows.

- 1. Modified sequencing algorithm is an non destructive in nature all the values of the job machine values are not deleted or removed.
- 2. Modified sequencing algorithm can be scaled to any size of job machine matrix.
- 3. Higher optimization results can be obtained when compared to the conventional algorithms.
- 4. The time taken to calculate or compute the solution is lesser when compared to conventional algorithm.
- 5. Faster iteration convergences along with the estimation of the convergence can be made.
- 6. Modified sequencing algorithm is linear in nature when coded and takes lower memory.
- 7. Operational logic is simple in nature.

## **4 NUMERICAL EXAMPLES.**

The following are the comparisons between the modified sequencing algorithm and the conventional sequencing algorithm.

**Example 1**: for the 4x2 job machine matrix shown below which is solved by conventional Johnson algorithm.

Job	Machine 1	Machine 2
А	3	4
В	5	2
С	2	3
D	6	2

**Sol:** check the smallest time between columns for the minimum values and placed them in front or end by the place of the value.

Job	Machine 1	Machine 2
А	3	4
В	5	2
С	2	3
D	6	2

Job	Machine 1	Machine 2
А	3	4
P	5	<u> </u>
D	5	2
С	2	3
D	6	2

Smallest value out of all cells is 2 which is in machine 2 so placed at end.

	D	,
		_

Job	Machine 1	Machine 2
А	3	4
B		2
-C-	2	
D	6	2

Smallest value out of all cells is 2 which is in the machine 1 so place the job in the front C

Job	Machine 1	Machine 2
А	3	4
<u>—B</u>		2
- <del>C</del> -	2	3
-Ð-	6	2

Smallest value out of all cells is 2 which is in the machine 2 so place the job in the end.

C D BThe optimum sequence is as above and the total time estimation is as follows.

Job	Mach	nine 1	Machine 2			
С	0	2	2	5		
А	2	5	5	9		
D	5	11	11	13		
В	11	16	16	18		

The time elapsed estimation is 18.

For the same problem application of the Modified sequencing algorithm which gives the following results.

Job	Machine 1	Machine 2						
А	3	4						
В	5	2						
С	2	3						
D	6	2						

Job	Machine	Machine	Minimum	Modified job
300	1	2	row values	sequence
Α	3	4	3	А
В	5	2	2	С
С	2	3	2	D
D	6	2	2	В

The optimum sequences of the jobs are as follows A C D B

		nalysis									Jobs	Machine 1	Machine 2	e Machii 3	ne Machir 4	ne Machin 5
		of <i>m</i> is gence		count	is <i>m-</i> '	1										
= 2-1=		Genee		count	10 111	1					A	8	6	2	3	9
		er of ic	ob int	er cha	nge is	n whi	ch is e	aual f	to 4.		В	6	6	4	5	10
				fficien						above	Č	5	4	5	6	8
		rified.			. 10 //						D	8	3	3	2	6
			able i	it is obs	serve	d that t	he op	timur	n sequ	uence		0	5	0	<u> </u>	0
				ods ar							Sol)					
				ne estir							,	ze of the jo	b machine	matrix is	4x5.	
		]	[ob	Machi	ine 1	Mach	nine 2	_			The co	nvergence	limit is <i>m</i>	-1 = 5-1=4		
		_	A	0	3	3	7	_				number of	iterations	are 4.		
			С	3	5	7	10				Iterati	on :1				
		1	D	5	11	11	13	-			ialaa	Machine	Machine	e Minin	num row	Modified jol
		]	В	11	16	16	18				jobs	1	2	va	alues	values
The v	alue	obtai	ned 1	here is	optir	nal the	n tha	t of t	he Iol	nson						
				lies th							А	8	6		3	D
son al			unere	nes u		ai vein	icatio	11 0010	ii uie	J0111-	В	6	6		4	С
	0		the	4x5 jo	h ma	chine	matri	v sho	wn ł	pelow	С	5	4		6	В
	-			-							D	8	3		6	А
winci	1 15 50	nveu b	y coi	nvectio	1101 01	goriui	III 15 a	5 10110	JW5.		Iterati	on :2				
	Mac	chine	Ma	chine	Ma	chine	Mac	hine	Ma	chine	Modi			Machine	Minimum	
lobs		1	1010	2		3		4	Ivita	5	job va	lues row	values	3	row value	s job value
А		8		6	-	2		3		9	D		3	2	3	А
В		6		6		4		5		10	C		4	4	5	В
С		5		4		5		6		8	В		6	5	4	C
D	:	8		3		3		2		6	A	2	6	3	2	D
Minir	num	proce	ssing	time	of N	lachine	e A≥	maxi	mum	pro-	Iteratio	n :3				
cessir	ng tirr	ne of N	/lachi	ine (B,	C, D	). This	case	is fail	ed be	cause						
5≥(6,											Modi	fied Mi	nimum	Machine	Minimum	Modified
				; time							job va	lues row	v values	4	row values	s job value
	0	ne of M	lachi	ne (B, G	C, D).	this ca	se is	worki	ng be	cause						
6≥(6 <i>,</i>											A		3	3	4	С
								hine p	proble	m by	B		5	5	2	D
Converting the 5 Machine problem to 2 Machine problem by assuming the imaginary Machine's G, H.											C		4	6	5	B
		A+B+									D		2	2	3	A
		=B+C+				~					Iterati	on ·4	-	-	0	
Wher							,									
		chine	Ma	chine	Ma	chine		thine	Ma	chine						10.110
Wher	Mac	-		2		3		4		5	Modi			Machine	Minimum	
Wher <u>The fo</u>		1		out	in	out	in	out	In	out	job va	lues row	v values	5	row values	s job value
Wher <u>The fo</u>	Mac	out	in	0			1(	19	19	28						
Wher <u>The fo</u>	Mac in	out			14	16	16			-	С		4	9	4	С
Wher <u>The fo</u> lobs 1	Mac in 0	out 8	8	14	14	16	16		•	07						
Wher <u>The fo</u> Jobs	Mac in	out			14 18	16 23	16 23	29	29	37	D		2	10	2	D
Wher <u>The fo</u> Jobs 1 3	Mac in 0 8	out 8	8 14	14 19		23	23						2 5		2 5	В
Wher <u>The fo</u> [obs 1 3 2	Mac in 0 8 13	out 8 13 19	8 14 19	14 19 25	18 25	23 29	23 29	29 34	37	47	D B A		2 5 3	10 8 6	2	
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for the 4x5 job machine matrix shown below which is solved T by modified sequencing algorithm.

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Job	Machine		chine Machine			chine	Ma	chine	Machine	
JOD	1		2		3		4		5	
	in	out	in	out	in	out	in	out	In	out
С	0	5	5	9	9	14	14	20	20	28
D	5	13	13	16	16	19	19	22	28	34
В	13	19	19	25	25	29	29	34	34	44
А	19	27	27	33	33	35	35	38	44	53

# CONCUSION

The observations proved that the application of the modified sequencing algorithm to n x 2 Job machine matrix and Johnson algorithm gives the same optimal solution, but application of the modified sequencing algorithm to  $n \ge 3$ ,  $n \ge 4$ ,  $n \ge m$  job machine matrix also give the same solution that of the convectional algorithm,but it should be the checking criteria for getting higher optimal solution along with the conventional algorithms because both cases fail in either way because some times in some problems they both may fail. For the n  $\ge 2$  problem it is not required to go for the next process if the optimum sequence for the both the methods are same.

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